

DASC 5333 Database Systems for Data Science
CSCI 4333 Design of Database Systems
Fall 2024
Suggested Solution for Homework #7

[1] See h7q1_sol.pdf.

[2] Proof of $F = \{A \rightarrow B, AB \rightarrow C, CD \rightarrow E, AE \rightarrow F\} \mid - AD \rightarrow F$

- [1] $A \rightarrow B$ (given)
- [2] $AB \rightarrow C$ (given)
- [3] $AA \rightarrow C$ (pseudo-transitivity on [1] and [2])
- [4] $A \rightarrow C$ (simplification of [3])
- [5] $CD \rightarrow E$ (given)
- [6] $AD \rightarrow E$ (pseudo-transitivity on [4] and [5])
- [7] $AE \rightarrow F$ (given)
- [8] $ADA \rightarrow F$ (pseudo-transitivity on [6] and [7])
- [9] $AD \rightarrow F$ (simplification of [8])

[3]

[a] $R(A,B,C,D) \{C \rightarrow D, AC \rightarrow B\}$

CK: [1] AC

Canonical Cover (optional): $\{C \rightarrow D, AC \rightarrow B\}$; as is.

Highest NF: 1NF

Reason: $C \rightarrow D$ violates 2NF since C is a proper subset of a CK, and D is non-prime.

[b] $R(A,B,C,D) \{C \rightarrow AD, D \rightarrow AB, A \rightarrow B, B \rightarrow A\}$

CK: [1] C

Canonical Cover (optional): $\{C \rightarrow D, D \rightarrow A, A \rightarrow B, B \rightarrow A\}$

Highest NF: 2NF

Reason: The FD $D \rightarrow A, A \rightarrow B,$ and $B \rightarrow A$ violate 3NF. The LHS are not superkeys and the RHS are non-prime attributes.

[c] $R(A,B,C,D) R(A,B,C,D) \{C \rightarrow BD, BD \rightarrow AC\}$

CK: [1] C, [2] BD

Canonical Cover (optional): $\{C \rightarrow BD, BD \rightarrow AC\}$; as is.

Highest NF: BCNF

Reason: The LHS of all non-trivial FD are superkeys.

[d] $R(A,B,C,D) \{C \rightarrow B, B \rightarrow C, BD \rightarrow A, CD \rightarrow A\}$

CK:[1] BD, [2] CD

Canonical Cover (optional): {C→B, B→C, BD→A}

Highest NF: 3NF

Reason: The FD C→B and B→C violate BCNF as B and C are both not superkeys.

[4] For F = {P→Q, R→S, PQ→ST, U→RS, UP→S}

(a)

P⁺=PQST

Q⁺= Q

R⁺ = RS

S⁺ = S

T⁺ = T

U⁺ = UPSR

(b) CK: [1] UP

L/NR: P, U

M: Q, R

R: S, T

(c) Prime: U,P; non-prime: Q, R, S, T

(d) Canonical Cover: {P→QST, R→S, U→R}

there are many, for examples, {RS→ Q, Q→RSTU, S→ P, U→Q}

(e) Highest NF: 1NF. S→P violates the 2NF.

(f) The following lossless decomposition results in all relations in BCNF.

R1(P,Q,R,S) {P→QST}

R2(R,S) {R→S}

R3(U,R) {U→R}

R4(P,U) {}

[5 For Employee(Eid, EFName, ELName, EPhone, DeptId, DeptName, PositionId, RateScale)

[a] Functional Dependencies (as in a canonical cover):

1. Eid → EId, EFName, ELName, PositionId
2. DeptId → DeptName
3. DeptName → DeptId (likely)
4. PositionId → RateScale

[b] The CKs are

1. Eid, DeptId, EPhone
2. Eid, DeptName, EPhone

[c] Thus, Highest normal form is 1NF as Eid → Eid, EFName, ELName, PositionId violate 2NF, for example.

[d] Decomposition:

Employee(Eid, EFName, ELName, PositionId)

1. FD: [1] Eid → Eid, EFName, ELName, PositionId
2. FK: [1] PositionId references Position(PositionId)
3. Highest NF: BCNF

EmployeePhone(Eid, EPhone)

1. FD: {}
2. FK: [1] Eid references Employee(Eid)
3. Highest Normal Form: BCNF

Department(DeptId, DeptName)

1. FD: [1] DeptId → DeptName, [2] DeptName → DeptId
2. FK: {}
3. Highest NF: BCNF

WorkFor(DeptId, Eid)

1. FD: {}
2. FK: [1] DeptId references Department(DeptId), [2] Eid references Employee(Eid)
3. Highest NF: BCNF

Position(PositionId, PayScale):

1. FD: [1] PositionId → PayScale
2. FK: {}
3. Highest NF: BCNF

[6] Minimum: 4, when there is only one candidate key: ABC.

Maximum: 28 when there are three candidate keys: [1] A, [2] B and [3] C.