

**DASC 5333 Database Systems for Data Science**  
**CSCI 4333 Design of Database Systems**  
**Spring 2025**  
**Suggested Solution for Homework #7**

[1] See h7q1\_sol.pdf.

[2]  $F = \{C \rightarrow AD, AC \rightarrow E, BE \rightarrow F, F \rightarrow A\}$  | - prove  $BC \rightarrow F$

Proof.

- [1]  $C \rightarrow AD$  (given)
- [2]  $C \rightarrow A$  (decomposition of [1])
- [3]  $AC \rightarrow E$  (given)
- [4]  $CC \rightarrow E$  (pseudo-transitivity rule on [2] and [3])
- [5]  $C \rightarrow E$  (simplification of [4])
- [6]  $BE \rightarrow F$  (given)
- [7]  $CB \rightarrow F$  (pseudo-transitivity rule on [5] and [6])

[3]

[a]  $R(A,B,C,D) \{B \rightarrow AD, AD \rightarrow C\}$

CK:[1] B

Canonical Cover (optional):  $\{B \rightarrow AD, AD \rightarrow C\}$ ; as is.

Highest NF: 2NF

Reason:  $AD \rightarrow C$  violates 3NF since AD is not a superkey and D is non-prime.

[b]  $R(A,B,C,D) \{AB \rightarrow D, BD \rightarrow C\}$

CK:[1] AB

Canonical Cover (optional):  $\{AB \rightarrow D, BD \rightarrow C\}$ ; as is.

Highest NF: 2NF

Reason: The FD  $BD \rightarrow C$  violate 3NF. BD is not a superkey and C is a non-prime attributes.

[c]  $R(A,B,C,D) \{C \rightarrow ABD, AD \rightarrow BC\}$

CK: [1] C, [2] AD

Canonical Cover (optional):  $\{C \rightarrow ABD, AD \rightarrow C\}$ .

Highest NF: BCNF

Reason: The LHS of all non-trivial FD are superkeys.

[d]  $R(A,B,C,D) \{A \rightarrow B, B \rightarrow A, AC \rightarrow D, AD \rightarrow C\}$

CK:[1] AC, [2] BC, [3] AD, [4] BD

Canonical Cover (optional):

Highest NF: 3NF

Reason: In 3NF as all attributes are prime. However,  $A \rightarrow B$  and  $B \rightarrow A$  violate BCNF.

[4] Consider  $F = \{A \rightarrow C, BD \rightarrow E, D \rightarrow BC, E \rightarrow B, BE \rightarrow C\}$

(a)  $A \neq AC, B \neq B, C \neq C, D \neq BCDE, E \neq BCE$

(b) The candidate key is AD

(c) Prime: A, D, non-prime: B, C, E

(d)  $\{A \rightarrow C, D \rightarrow E, E \rightarrow BC\}$

(e) 1NF since  $A \rightarrow C$  violates 2NF: C is non-prime and A is a proper subset of a CK.

(f) Yes, the decomposition:

1.  $R_1(A, C) \{A \rightarrow C\}$
2.  $R_2(D, E) \{D \rightarrow E\}$
3.  $R_3(B, C, E) \{E \rightarrow BC\}$
4.  $R_4(A, D) \{\}$

[5] (a) FD:

FD1:  $studId \rightarrow stuName$

FD2:  $studId, semester \rightarrow dormName, dormRoom, mealPlan$

FD3:  $dormName, dormRoom, semester \rightarrow dormFee$

FD4:  $mealPlan, semester \rightarrow mealFee$

(b) CK:  $\{studId, semester\}$

(c) 1NF. FD1 violates 2NF. FD3 and FD4 violate 3NF.

(d) Decomposition into BCNF component relations:

Student( $studId, stuName, ..$ )  $\{studId \rightarrow stuName\}$

DormFee( $dormName, dormRoom, semester, dormFee$ )  $\{dormName, dormRoom, semester \rightarrow dormFee\}$

MealFee( $mealPlan, semester, mealFee$ )  $\{mealPlan, semester \rightarrow mealFee\}$

StudentSemester( $studId, semester, dormName, dormRoom, mealPlan$ )  $\{studId, semester \rightarrow dormName, dormRoom, mealPlan\}$

[6] Minimum: 9, e.g., when ACDE is the second candidate key.

Maximum: 20, e.g., when C is the second candidate key.