

DASC 5333 Database Systems for Data Science
CSCI 4333 Design of Database Systems
Spring 2025
Suggested Solution for Homework #7

[1] See h7q1_sol.pdf.

[2] $F = \{C \rightarrow AD, AC \rightarrow E, BE \rightarrow F, F \rightarrow A\}$ | - prove $BC \rightarrow F$

Proof.

- [1] $C \rightarrow AD$ (given)
- [2] $C \rightarrow A$ (decomposition of [1])
- [3] $AC \rightarrow E$ (given)
- [4] $CC \rightarrow E$ (pseudo-transitivity rule on [2] and [3])
- [5] $C \rightarrow E$ (simplification of [4])
- [6] $BE \rightarrow F$ (given)
- [7] $CB \rightarrow F$ (pseudo-transitivity rule on [5] and [6])

[3]

[a] $R(A,B,C,D) \{B \rightarrow AD, AD \rightarrow C\}$

CK: [1] B

Canonical Cover (optional): $\{B \rightarrow AD, AD \rightarrow C\}$; as is.

Highest NF: 2NF

Reason: $AD \rightarrow C$ violates 3NF since AD is not a superkey and D is non-prime.

[b] $R(A,B,C,D) \{AB \rightarrow D, BD \rightarrow C\}$

CK: [1] AB

Canonical Cover (optional): $\{AB \rightarrow D, BD \rightarrow C\}$; as is.

Highest NF: 2NF

Reason: The FD $BD \rightarrow C$ violate 3NF. BD is not a superkey and C is a non-prime attributes.

[c] $R(A,B,C,D) R(A,B,C,D) \{C \rightarrow ABD, AD \rightarrow BC\}$

CK: [1] C, [2] AD

Canonical Cover (optional): $\{C \rightarrow ABD, AD \rightarrow C\}$.

Highest NF: BCNF

Reason: The LHS of all non-trivial FD are superkeys.

[d] $R(A,B,C,D) \{A \rightarrow B, B \rightarrow A, AC \rightarrow D, AD \rightarrow C\}$

CK: [1] AC, [2] BC, [3] AD, [4] BD

Canonical Cover (optional):

Highest NF: 3NF

Reason: In 3NF as all attributes are prime. However, $A \rightarrow B$ and $B \rightarrow A$ violate BCNF.

[4] Consider $F = \{A \rightarrow C, BD \rightarrow E, D \rightarrow BC, E \rightarrow B, BE \rightarrow C\}$

(a) $A \rightarrow AC, B \rightarrow B, C \rightarrow C, D \rightarrow BCDE, E \rightarrow BCE$

(b) The candidate key is AD

(c) Prime: A, D, non-prime: B, C, E

(d) $\{A \rightarrow C, D \rightarrow E, E \rightarrow BC\}$

(e) 1NF since $A \rightarrow C$ violates 2NF: C is non-prime and A is a proper subset of a CK.

(f) Yes, the decomposition:

1. $R_1(A, C) \{A \rightarrow C\}$
2. $R_2(D, E) \{D \rightarrow E\}$
3. $R_3(B, C, E) \{E \rightarrow BC\}$
4. $R_4(A, D) \{\}$

[5] (a) FD:

FD1: $stuld \rightarrow stulname$

FD2: $studId, semester \rightarrow dormName, dormRoom, mealPlan$

FD3: $dormName, dormRoom, semester \rightarrow dormFee$

FD4: $mealPlan, semester \rightarrow mealFee$

(b) CK: $\{stuld, semester\}$

(c) 1NF. FD1 violates 2NF. FD3 and FD4 violate 3NF.

(d) Decomposition into BCNF component relations:

Student($stuld, stulname, ..$) $\{ stuld \rightarrow stulname \}$

DormFee($dormName, dormRoom, semester, dormFee$) $\{ dormName, dormRoom, semester \rightarrow dormFee \}$

MealFee($mealPlan, semester, mealFee$) $\{ mealPlan, semester \rightarrow mealFee \}$

StudentSemester($studId, semester, dormName, dormRoom, mealPlan$) $\{ studId, semester \rightarrow dormName, dormRoom, mealPlan \}$

[6] Minimum: 9, e.g., when ACDE is the second candidate key.

Maximum: 20, e.g., when C is the second candidate key.